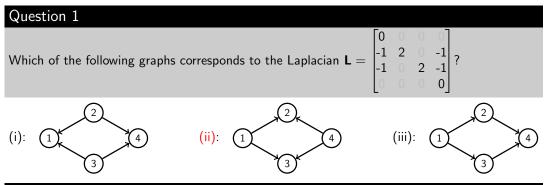
Advanced Systems Theory Lecture 11: Multiagent systems: Consensus

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Recap graph theory



Question 2

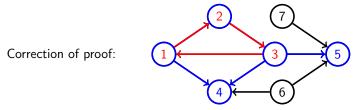
Which of the following statements is correct for the above graph?

- (i) $\lambda_2 = 0$ because the graph contains an (undirected) cycle
- (ii) $\lambda_2 = 0$ because rank(L) = 2
- (iii) $\lambda_2 > 0$ because it consists of two overlapping spanning trees
- (iv) $\lambda_2 > 0$ because it is connected

Proof of "no spanning tree $\implies \lambda_2 = 0$ " revisited

Question 3

Assume G has a spanning tree. Is it true that then there is a root of the spanning tree with (in-)degree zero?



Step 1: Choose a maximal tree $(\mathcal{V}_{\mathcal{T}}, \mathcal{E}_{\mathcal{T}})$

Step 2: Choose a maximal root set $\mathcal{V}_R \subseteq \mathcal{V}_T$ and relabel nodes

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The consensus problem for simple integrator dynamics

The consensus problem: setup

Given

N agents with simple scalar integrator dynamics:

$$\dot{\mathbf{x}}_i = \mathbf{u}_i, \quad \mathbf{x}_i(0) = \mathbf{x}_i^0 \in \mathbb{R}$$

Communication graph G = (V, E) describing which information the agents can access from their neighbors

Objective

Design coupling laws $(x_i, x_{\mathcal{N}_i}) \mapsto u_i(x_i, x_{\mathcal{N}_i})$ such that

 $\forall i \in \mathcal{V}: \quad x_i(t) \to \overline{x} \text{ as } t \to \infty$

In particular, $x_i(t) - x_j(t) \rightarrow 0$ as $t \rightarrow \infty$.

Attention: \overline{x} not known by agents a priori

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The consensus protocol for $\dot{x}_i = u_i$

Feasible but undesired choice of feedback

 $u_i = -x_i$ leads to $\dot{x}_i = -x_i \quad \rightsquigarrow \quad x_i(t) \to 0$ exponentially

Additional requirement

If for some $t \ge 0$ $x_i(t) = x_j(t)$ for all $j \in \mathcal{N}_i$ (including $\mathcal{N}_i = \emptyset$) then $u_i(t) = 0$.

Simplest case: Agent i has one neighbor j

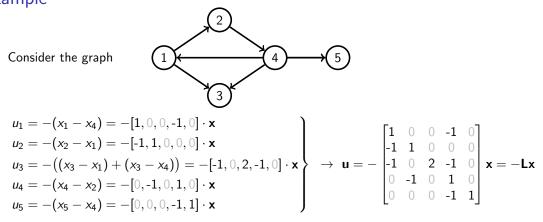
$$egin{aligned} & x_i(t) < x_j(t) & \stackrel{!}{ o} & \dot{x}_i(t) > 0 \ \leftrightarrow \ u_i(t) > 0 \ x_i(t) > x_j(t) & \stackrel{!}{ o} & \dot{x}_i(t) < 0 \ \leftrightarrow \ u_i(t) < 0 \ \end{aligned}
ightarrow egin{aligned} & u_i(t) > 0 \ x_i(t) > x_j(t) & \stackrel{!}{ o} & \dot{x}_i(t) < 0 \ \leftrightarrow \ u_i(t) < 0 \ \end{aligned}$$

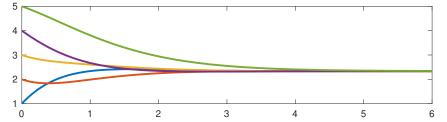
More then one neighbor \rightarrow Diffusive coupling

$$u_i(t) = -\sum_{j\in\mathcal{N}_i} (x_i(t) - x_j(t))$$

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Example

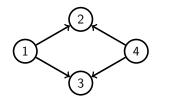




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Another example

Consider now the graph

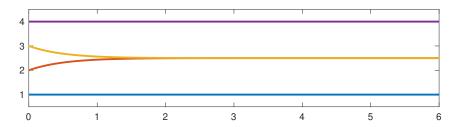


$$u_1 = 0$$

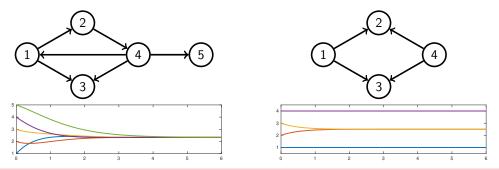
$$u_2 = -((x_2 - x_1) + (x_2 - x_4))$$

$$u_3 = -((x_3 - x_1) + (x_3 - x_4))$$

$$u_4 = 0$$



Consensus: When and what?



Key questions

- 1. Under which conditions does consensus occur?
- 2. What is the consensus value?

Brain storming (use chat for answers)

Which property of the graph seems crucial for consensus? On what does the consensus value probably depend? Can you say something about the location of the consensus value?

Consensus: Necessary and sufficient condition

Connectivity

- ▶ Clearly, if \mathcal{G} has two (or more) completely disconnected components \rightsquigarrow no consensus
- However, strong connectivity doesn't seem to be necessary (cf. first example)
- Spanning tree necessary? ... Yes!

Desired consensus value is $\mathbf{x}(\infty) = \overline{x} \cdot \mathbf{1}$, where $\mathbf{1} = (1, 1, \dots, 1)^\top \in \mathbb{R}^n$

Consensus feasible under diffusive coupling?

- ▶ Is consensus value equilibrium of closed loop $\dot{x} = -Lx$? Yes! L1 = 0
- Consensus only possible for all x(0) if ker L = span{1}
 - $\ker \mathbf{L} = \operatorname{span}\{\mathbf{1}\} \iff \dim \ker \mathbf{L} = \mathbf{1} \iff \lambda_2(\mathbf{L}) \neq \mathbf{0} \iff \mathcal{G} \text{ has spanning tree}$

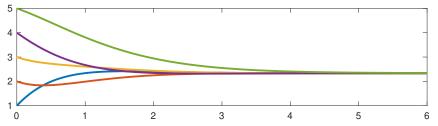
Theorem

Diffusive coupling leads to consensus $\iff \mathcal{G}$ has a spanning tree

Proof ...

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The consensus value for $\dot{\mathbf{x}} = -\mathbf{L}\mathbf{x}$



Lemma (Proof: Exercise)

For $t_0 \ge 0$ let $x_{\min}(t_0) := \min_{j \in \mathcal{V}} x_j(t_0)$ and $x_{\max}(t_0) := \max_{j \in \mathcal{V}} x_j(t_0)$, then $\forall t \ge t_0 \ \forall i \in \mathcal{V} : \quad x_i(t) \in [x_{\min}(t_0), x_{\max}(t_0)]$

Key idea to find consensus value: Try to find invariant

Is there a linear combination $\mathbf{w}^{\top}\mathbf{x}(t)$ which remains constant (i.e. which is an invariant)?

Brain storming (use chat for answers)

Which vectors $\mathbf{w} \in \mathbb{R}^N$ lead to an invariant $\mathbf{w}^\top \mathbf{x}(t)$?

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Calculation of consensus value

Lemma

$$\mathbf{w}^{\top}\mathbf{x}(t)$$
 is invariant for $\dot{\mathbf{x}} = -\mathbf{L}\mathbf{x} \iff \mathbf{w} \in \ker \mathbf{L}^{\top}$

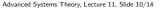
Furthermore, if rank $\mathbf{L} = N - 1$ then there exists unique $\widehat{\mathbf{w}} \in \ker \mathbf{L}^{\top}$ with $\widehat{\mathbf{w}}^{\top} \mathbf{1} = \mathbf{1}$, in fact

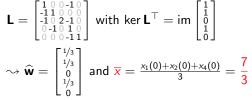
$$\widehat{\mathbf{w}} = rac{\mathbf{w}}{\mathbf{w}^{ op}} \mathbf{1}$$
 and $\widehat{\mathbf{w}} \geq 0$

Corollary

 \mathcal{G} has spanning tree $\implies \mathbf{x}(t) \to \overline{\mathbf{x}} \cdot \mathbf{1}$ with $\overline{\mathbf{x}} = \widehat{\mathbf{w}}^\top \mathbf{x}(0)$ Proof ...

 $\mathbf{L} = \begin{bmatrix} 1 & 0 & 0 & 1 & 0 \\ -1 & 1 & 0 & 0 & 0 \\ -1 & 0 & 2 & 1 & 0 \\ 0 & 0 & 0 & -1 & 1 \end{bmatrix} \text{ with}$





Interpretation of consensus value

Lemma (without proof)

The i-th component of $\widehat{\mathbf{w}}$ is non-zero \iff all nodes can be reached from node i

Corollary

 \overline{x} only depends on initial values of agents which are roots of spanning trees

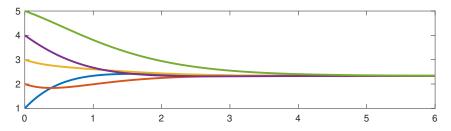
Corollary: Undirected case

Let \mathcal{G} be undirected and connected, then \overline{x} is the average of initial values, i.e.

$$\overline{x} = rac{1}{N}\sum_{i=1}^N x_i(0)$$

In particular, it is independent of the underlying graph structure!

Convergence rate



Question

How fast do the agents converge towards the consensus value?

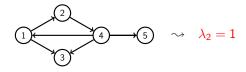
Theorem

$$|x_i(t) - \overline{x}| \leq c_{\mathbf{x}(0)} e^{-\operatorname{\mathsf{Re}}(\lambda_2)t}$$

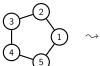
Reminder: For any Hurwitz matrix M we have $||e^{Mt}|| \le ce^{-\lambda_{\min}t}$ for some c > 0 and $\lambda_{\min} = \in \{ |\operatorname{Re}(\lambda)| \mid \lambda \text{ is eigenvalue of } M \}$ Proof: ...

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Convergence rate for specific graphs







 $\lambda_2 pprox 0.382$

$$\begin{array}{c} 3 \\ 3 \\ 4 \\ 4 \\ 5 \end{array} \right)^{2} \rightarrow \lambda_{2} = \lambda_{2}$$

$$\rightarrow \lambda_2 = 3$$

Summary

The consensus problem: The integrator dynamics case

- Agent dynamics: $\dot{x}_i = u_i$
- ▶ Goal: $x_i(t) \to \overline{x}$ for all $i \in \mathcal{V}$
- Diffusive coupling: $u_i = -\sum_{j \in \mathcal{N}_i} (x_i x_j)$
- Resulting overall dynamics: x = -Lx

Consensus results for simple multi-agent system

• Consensus $\iff \mathcal{G}$ has spanning tree (undirected case: is connected)

- $\mathbf{x} = \widehat{\mathbf{w}}^{\top} \mathbf{x}(0)$ where $\widehat{\mathbf{w}}^{\top} L = 0$ and $\mathbf{w}^{\top} \mathbf{1} = 1$ (undirected case: average of initial values)
- Convergence rate proportional to $e^{-\operatorname{Re}(\lambda_2)t}$